

# Hidden Metric Spaces and Navigability of Complex Networks

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# ~~Science or engineering?~~

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- ~~# Network science vs. network engineering~~
  - ~~# Computer science vs. computer engineering~~
  - ~~# Study existing networks vs. designing new ones~~
  - # We cannot really *design* truly large-scale systems (e.g., Internet)
    - We can design their building blocks (e.g., IP)
    - But we cannot fully control their large-scale behavior
    - At their large scale, complex networks exhibit some emergent properties, which we can only observe: we cannot yet fully understand them, much less predict, much less control
  - # Let us study existing large-scale networks and try to use what we learn in designing new ones
    - Discover “nature-designed” efficient mechanisms that we can reuse (or respect) in our future designs
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# Internet

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- # Microscopic view (“designed constraints”)
  - IP/TCP, routing protocols
  - Routers
  - Per-ISP router-level topologies
- # Macroscopic view (“non-designed emergent properties”)
  - Global AS-level topology is a cumulative result of local, decentralized, and rather complex interactions between AS pairs
  - Surprisingly, in 1999, it was found to look completely differently than engineers and designers had thought
    - It is not a grid, tree, or classical random graph
    - It shares all the main features of topologies of other complex networks
      - scale-free (power-law) node degree distributions ( $P(k) \sim k^{-\gamma}$ ,  $\gamma \in [2,3]$ )
      - strong clustering (large numbers of 3-cycles)

# Problem

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- # “Designed parts” have to deal with “emergent properties”
  - For example, BGP has to route through the existing AS topology, which was not a part of BGP design

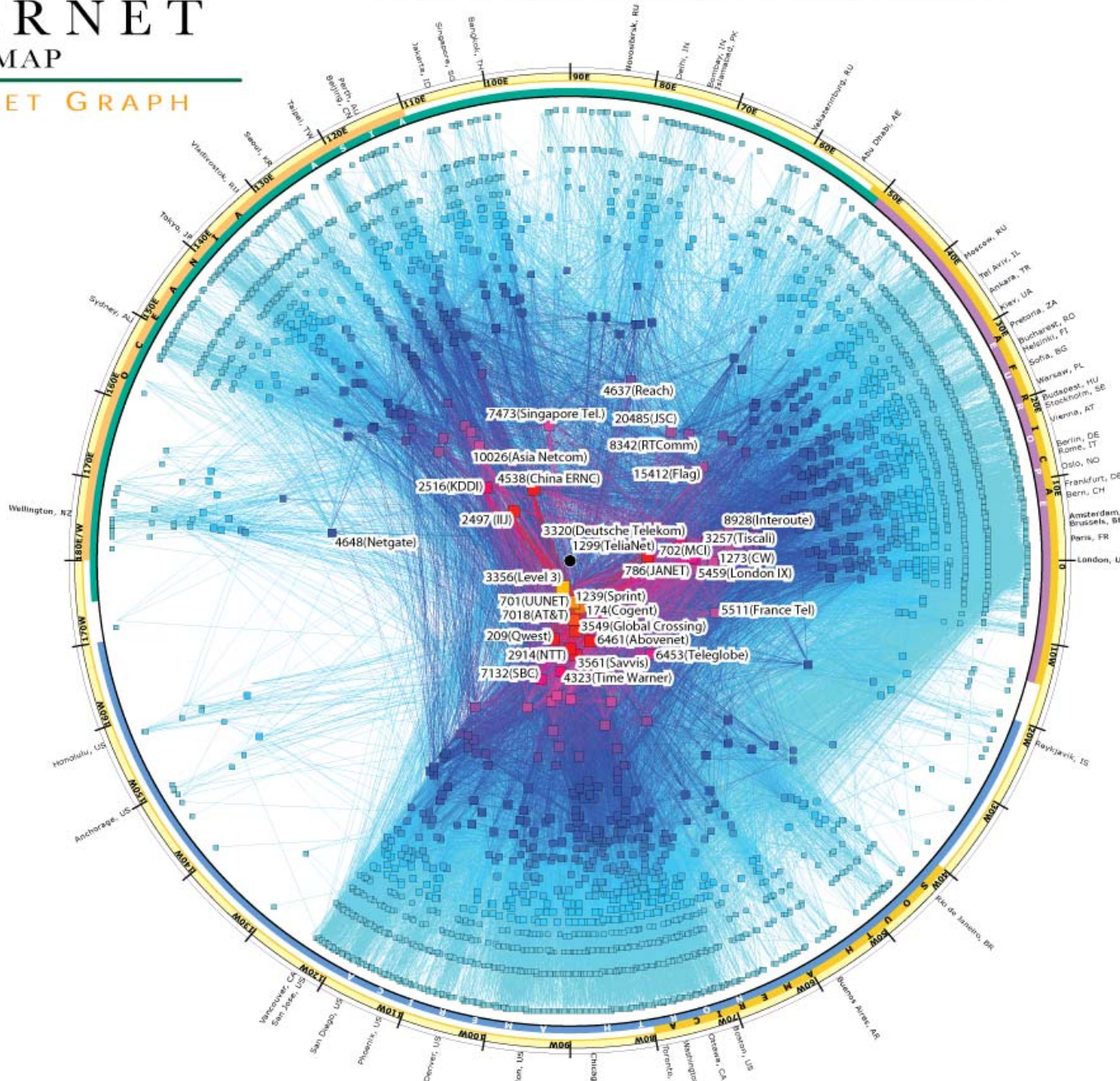
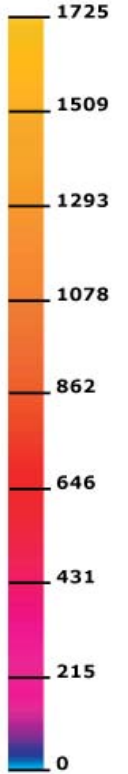


# IPv4 INTERNET TOPOLOGY MAP

## AS-level INTERNET GRAPH

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Peering:  
OutDegree



# Routing practice

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## # Global (DFZ) routing tables

- 300,000 prefix entries (and growing)
- 30,000 ASs (and growing)

## # Routing overhead/convergence

- BGP updates
  - 2 per second on average
  - 7000 per second peak rate
- Convergence after a single event can take up to tens of minutes

## # Problems with design?

- Yes and no
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# Routing theory

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- # There can be no routing algorithm with the number of messages per topology change scaling better than linearly with the network size in the worst case
- # Small-world networks are this worst case

CCR, v.37, n.3, 2007

- # *Is there any workaround?*
  - # *If topology updates/convergence is so expensive, then may be we can route without them, i.e., without global knowledge of the network topology?*
  - # *What about other existing networks?*
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# Navigability of complex networks

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- # In many (if not all) existing complex networks, nodes communicate without any global knowledge of network topologies; examples:
    - Social networks
    - Neural networks
    - Cell regulatory networks
  - # How is this possible???
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# Hidden metric space explanation

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- # All nodes exist in a metric space
  - # Distances in this space abstract node similarities
    - More similar nodes are closer in the space
  - # Network consists of links that exist with probability that decreases with the hidden distance
    - More similar/close nodes are more likely to be connected
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# Mathematical perspective: Graphs embedded in manifolds

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- # All nodes exist in “two places at once”:
  - graph
  - hidden metric space, e.g., a Riemannian manifold
- # There are two metric distances between each pair of nodes: observable and hidden:
  - hop length of the shortest path in the graph
  - distance in the hidden space

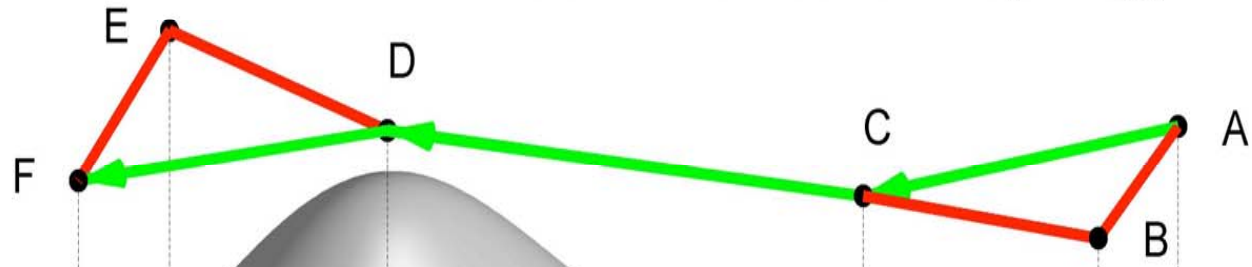


# Greedy routing (Kleinberg)

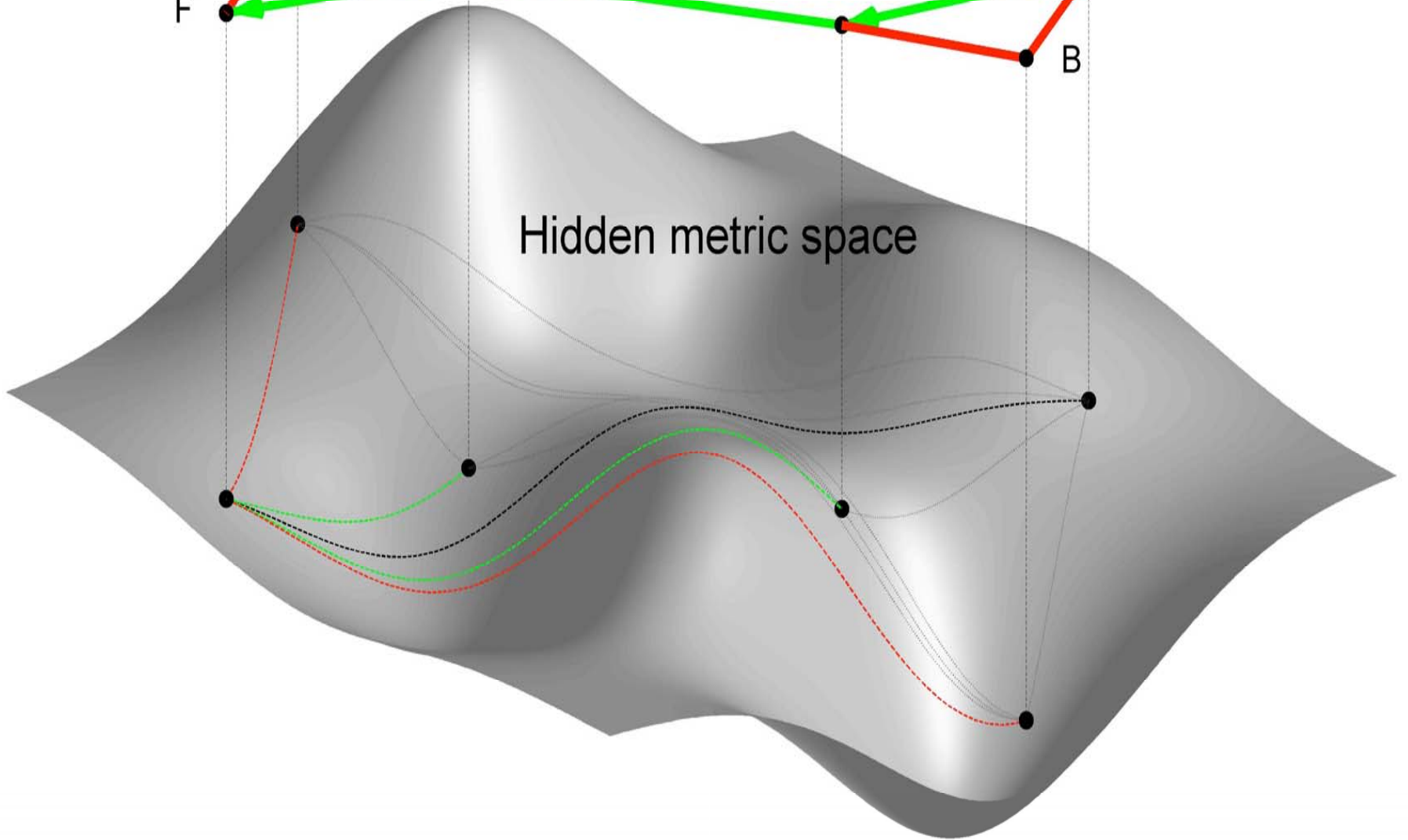
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- # To reach a destination, each node forwards information to the one of its neighbors that is closest to the destination in the hidden space

# Observable network topology



# Hidden metric space





# Result #1:

## Hidden metric space do exist

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- # Their existence appears as the only reasonable explanation of one peculiar property of the topology of real complex networks – self-similarity of clustering

## Result #2:

Complex network topologies are navigable

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- # Specific values of degree distribution and clustering observed in real complex networks correspond to the highest efficiency of greedy routing
- # Which implicitly suggests that complex networks do evolve to become navigable
- # Because if they did not, they would not be able to function



## Result #3:

# Successful greedy paths are shortest

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- # Regardless the structure of the hidden space, complex network topologies are such, that all successful greedy paths are asymptotically shortest
- # But: how many greedy paths are successful does depend on the hidden space geometry

# Result #4:

In hyperbolic geometry, all paths are successful

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- # Hyperbolic geometry is the geometry of trees; the volume of balls grows exponentially with their radii
  - # Greedy routing in complex networks, including the real AS Internet, embedded in hyperbolic spaces, is always successful and always follows shortest paths
  - # Even if some links are removed, emulating topology dynamics, greedy routing finds remaining paths if they exist, without recomputation of node coordinates
  - # The reason is the exceptional congruency between complex network topology and hyperbolic geometry
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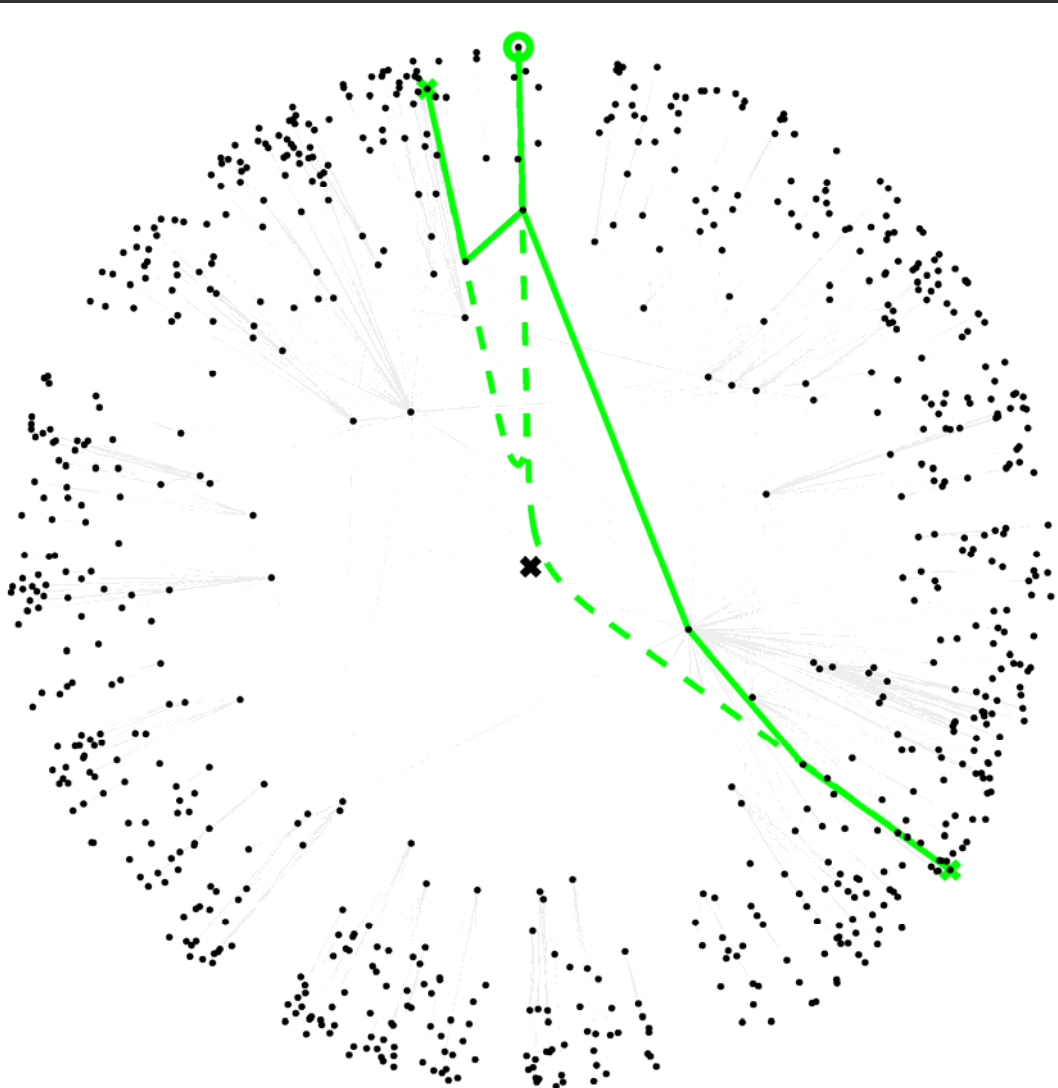
# Result #5:

## Emergence of topology from geometry

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- # The two main properties of complex network topology are direct consequences of the two main properties of hyperbolic geometry:
    - Scale-free degree distributions are a consequence of the exponential expansion of space in hyperbolic geometry
    - Strong clustering is a consequence of the fact that hyperbolic spaces are metric spaces
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# Shortest paths in scale-free graphs and hyperbolic spaces

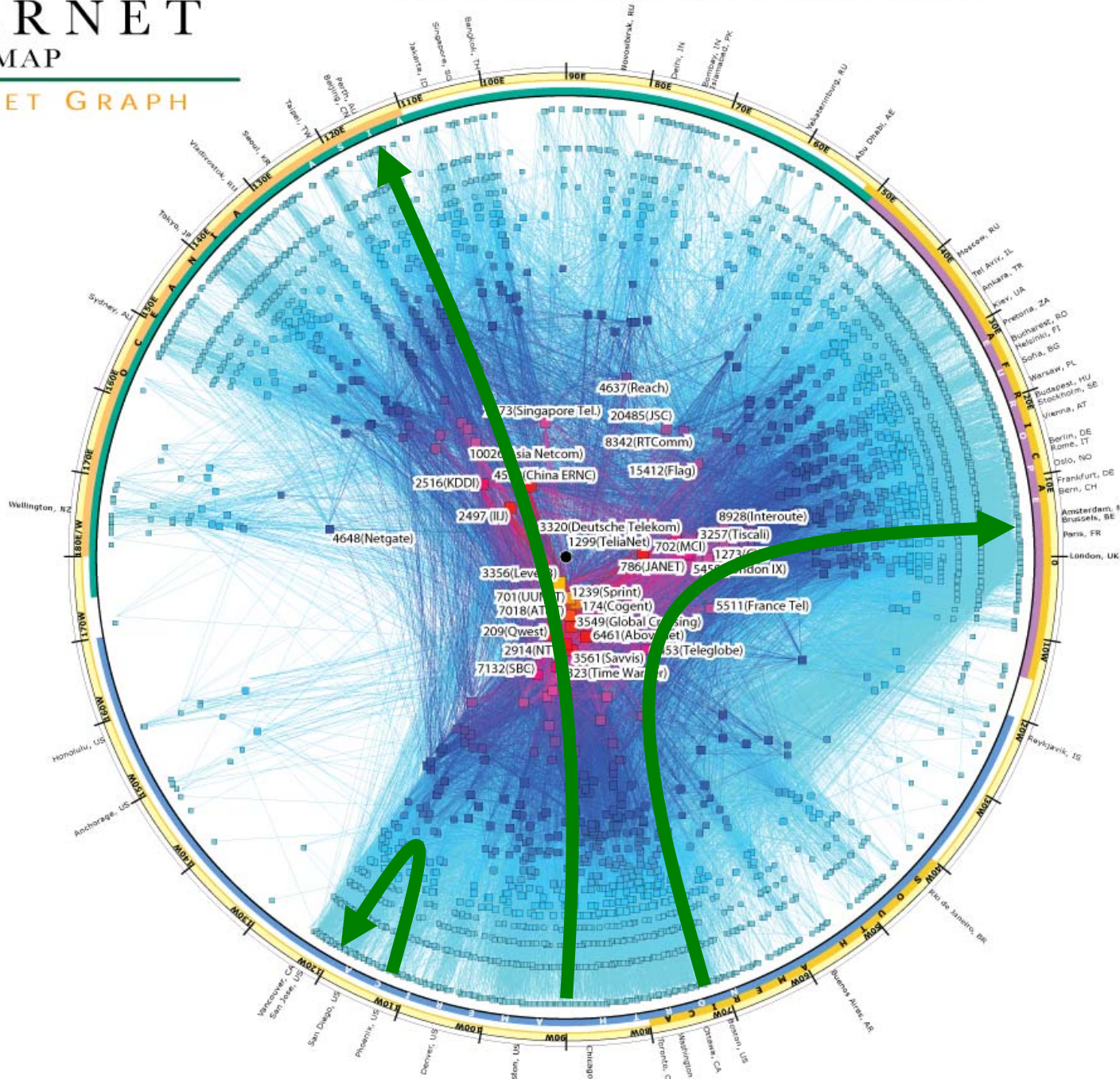
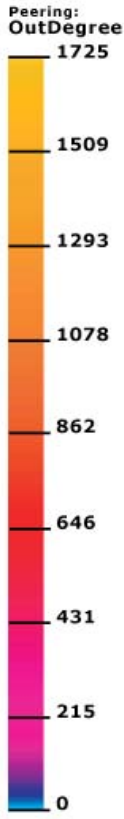




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## AS-level INTERNET GRAPH

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# In summary

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- # Complex network topologies are congruent with hidden hyperbolic geometries
  - Greedy paths follow shortest paths that approximately follow shortest hidden paths, i.e., geodesics in the hyperbolic space
    - Both topology and geometry are *tree-like*
- # This congruency is robust w.r.t. topology dynamics
  - There are many link/node-disjoint shortest paths between the same source and destination that satisfy the above property
    - *Strong clustering* (many by-passes) boosts up the path diversity
  - If some of shortest paths are damaged by link failures, many others remain available, and greedy routing still finds them



# Conclusion

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- # To efficiently route without topology knowledge, the topology should be both hierarchical (tree-like) and have high path diversity (not like a tree)
  - # Complex networks do borrow the best out of these two seemingly mutually-exclusive worlds
  - # Hidden hyperbolic geometry naturally explains how this balance is achieved
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# Applications

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- # Greedy routing mechanism in these settings may offer virtually infinitely scalable information dissemination (routing) strategies for future communication networks
    - Zero communication costs (no routing updates!)
    - Constant routing table sizes (coordinates in the space)
    - No stretch (all paths are shortest, stretch=1)
  - # Interdisciplinary applications
    - systems biology: brain and regulatory networks, cancer research, phylogenetic trees, protein folding, etc.
    - data mining and recommender systems
    - cognitive science
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