Segmentation of Internet Paths for Capacity Estimation

Matthew Luckie (PhD Student) Tony McGregor (Supervisor)

WAND Group, University of Waikato NLANR Measurement and Network Analysis

mjl@luckie.org.nz tonym@cs.waikato.ac.nz

Introduction

 We are talking about capacity estimation of individual hops in an end-to-end path

 If we had a method to segment a path into hops, what problems might it solve, and what might remain?

Current Techniques

opathchar - Jacobson, original TTL method

° clink and pchar

 Downey suggests statistical methods to detect convergence on a link capacity faster

What is hard?

 Capacity estimation of each hop in a path is hard to do with the TTL method

additive noise from prior links brings doubt

can only estimate the forward path

L2 store-and-forward devices interfere

• each hop is probed with a separate series of packets

A goal

 A method to segment a path into links could be useful

remove additive error

• measure capacity of each link in a long path

 Identified as future research by Jacobson in MSRI talk

Other goals

 Measure the capacity of hops on the reverse path

• Be kind to the network

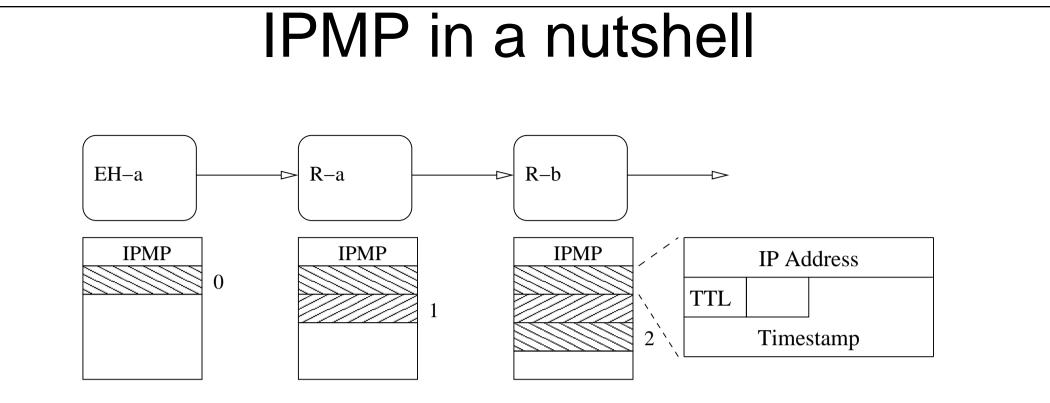
 Measure capacities of hops correctly when Layer-2 Store and Forward devices are in the path

A method

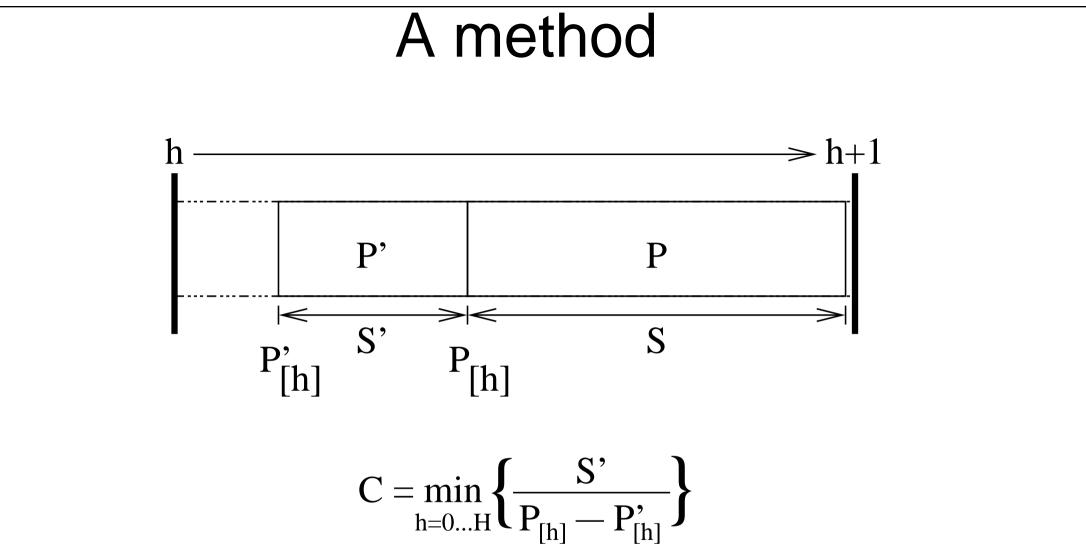
 Use timestamps inserted at each hop [modify the forwarding path, IPMP]

 Use the packet-tailgating technique [nettimer] to encourage a packet pair to queue together through a network

 (extremely simplified) Estimate the capacity of each hop using the time difference between the first packet and the second packet



A specially-marked packet passes through a network, collecting timestamps in "path records" as it goes



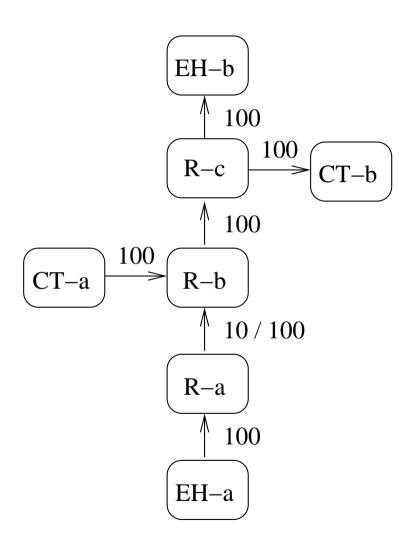
What does this buy us?

 the behaviour of each segment is confined to that segment unless the packets become separated [more on this shortly]

• measures the forward and reverse paths

 measures the capacity of hops with L2 store-and-forward devices correctly as it uses packet-pairs

Experiments



EH = Echo Host R = Router CT = Cross Traffic Host

Linux 2.4.20 + PPS Kit + IPMP Kernel Implementation + Crossover Cables

Arrows indicate direction from source to destination

R-a to R-b is forced 10mbps with a 10mbps 3COM hub in some experiments

An aside: ctft

Wrote cross-traffic-from-trace (ctft)

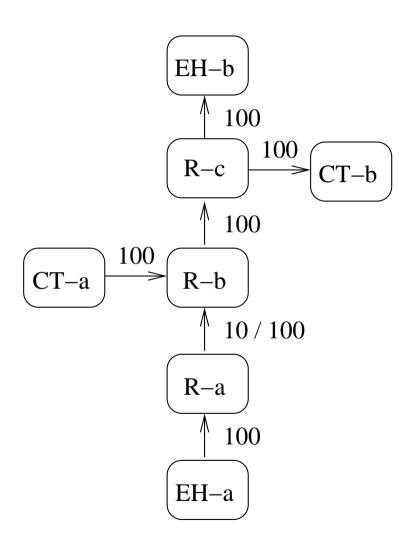
using an Auckland trace

given a probability of a packet of a given size

and the probability of a time-to-next packet

send combinations of size and delay randomly that fits the profile

Experiments



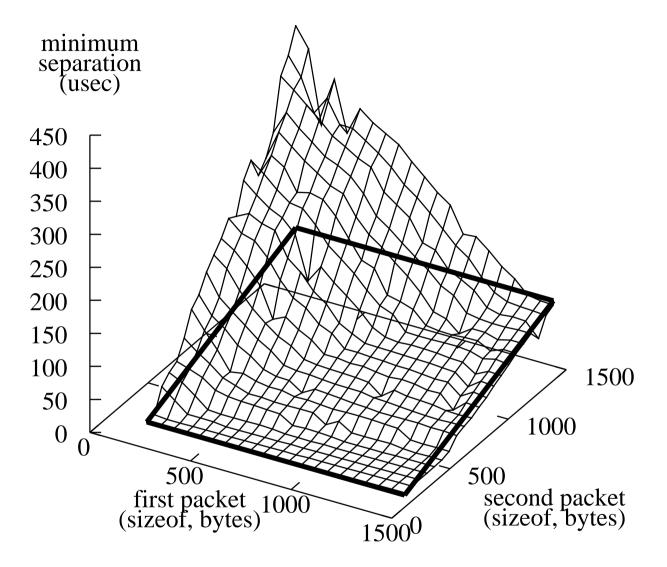
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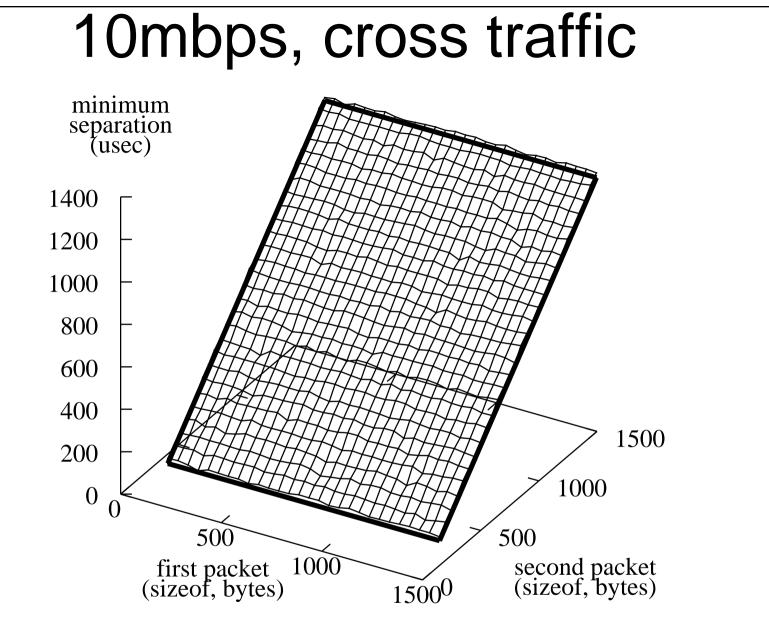
Linux 2.4.20 + PPS Kit + IPMP Kernel Implementation + Crossover Cables

Arrows indicate direction from source to destination

R-a to R-b is forced 10mbps with a 10mbps 3COM hub in some experiments

100mbps, no cross traffic, Hop 4





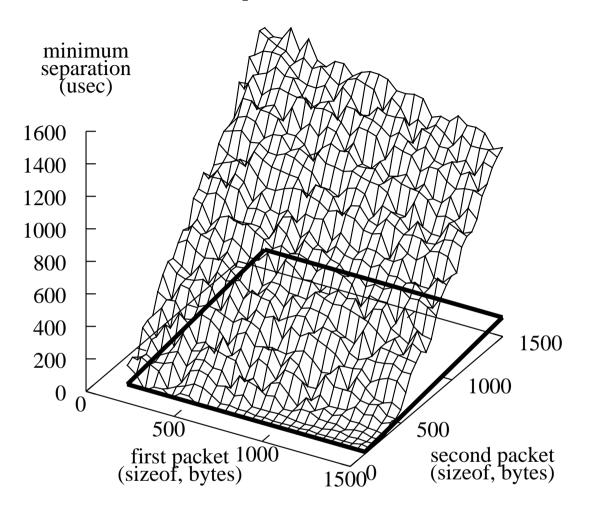
•What problems still remain?

 given an egress link n times faster than the ingress link to a router, then S' must be no larger than S/n if the pair is to remain together

 i.e. measuring a 100mbps hop immediately after a 10mbps hop is problematic; cross traffic can help here, though

Imited number of path records in a small packet

10->100mbps, cross traffic



Disadvantages

One of the protocol deployed

 not everywhere before it becomes useful for capacity estimation, though

 can estimate the capacity of the path through two points by looking for the minimum time for a pair through the points and the minimum separation of said pair

Future work

 work at estimating congestion / queue length at intermediate nodes

Olooking for Layer 2 capacities

cross traffic (CT) is likely to be either a 40 / 576 / 1500 byte packet

Iook for CT separating the pair (cross traffic) at various sizes and infer what the capacity of the link was that caused the separation

Conclusion

 Presented a method to segment a path into hops for capacity estimation

- Discussed the advantages ...
- $^{\rm O} \dots$ and caveats

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802.11b Network [an aside]

